

Finding Simplex Items in Data Streams

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Outline

PART 01 Background PART 02 Problem Statement PART 03 X-Sketch Design PART 04 Experimental Results PART 05 X-Sketch for ML PART 06 Conclusion

PART ONE Background

01

01/ Background

Frequency estimation is important

Finding frequent items in data streams have been well studied by the community .

Best solution: Sketch

Memory efficient and time efficient
 Small errors

01/Background

Patterns in which item frequencies present in a certain number of <u>consecutive</u> windows are also worth exploring.

Such patterns in p consecutive windows may be fitted by k-degree polynomials.



01/ Background

k=0:

1) stable cache lines

2) stable flows in network management



01/ Background

k=1:

speed up machine learning models detect DDoS attacks



01/Background

k=2:

1) monitor traffic in wireless networks



01/ Background- CMSketch



01/ Background- TowerSketch



O2 PART TWO

Problem Statement

02/ Problem Statement

k-simplex items

An item e is called k-simplex item from window wif and only if its frequencies in p consecutive windows $f_w, f_{w+1}, \cdots, f_{w+p-1}$ satisfy: (1) $f_{w+i} > 0, \forall i = 0, 1, \cdots, p-1;$ (2) There exists a k^{th} -degree polynomial $f(n) = a_0 + a_1 n + a_2 n^2 + \dots + a_k n^k = \sum_{j=0}^{n} a_j n^j,$ such that the mean squared error (MSE) ε satisfies $\varepsilon = \frac{1}{p} \sum_{i=0}^{p-1} (f(i) - f_{w+i})^2 \le T;$

where T is a predefined threshold.

02/ Problem Statement

simplex items and persistent items

number of windows items appear
 frequency is important
 consecutive windows



X-Sketch

01 Baseline Solution
02 An Optimization
03 Data Structure

03 X-Sketch - Baseline Solution

Part 1:

Record frequencies for last p windows using CMSketch

Part 2:

Record lasting time of simplex items using a set and a hash table.

03 X-Sketch - Optimization

The fitting error will not increase if we simply increase k.



Distinguish k-simplex items from (k-1)-simplex items: Threshold for |a_k|



03/X-Sketch - Data Structure

Stage 1:

Record frequencies for last s windows using TowerSketch.

Stage 2:

Record potential simplex items and their lasting time using a hash table.

03 X-Sketch - Data Structure

Stage 1:

Short-Term Filtering: quickly filter non-simplex items
 Potential: measure the probability of becoming simplex items

$$\Lambda = \frac{|a_k|}{\varepsilon + \Delta}$$

03/X-Sketch - Data Structure

Stage 2:

Weight Election: a reasonable replacement strategy

03/X-Sketch - Data Structure





PART FOUR Experimental Results

04 Experiments (Setup)

Datasets: CAIDA, MAWI, Data Center, Synthetic.

Metrics:

PR, RR, F1 Score, ARE, Throughput

04 Experiments (F1 Score)



04 Experiments (ARE)



04 Experiments (Throughput)





05/X-Sketch for ML

ML for predicting frequency:
1) massive training datasets
2) loops of training epochs
3) not all items are predictable

Solution: X-Sketch + ML



ML

X-Sketch+ML

05/X-Sketch for ML

TABLE III: Experiments on the Transactional Dataset.

	Model	Accuracy (%)	Running Time (s)
k = 0	X-Sketch (C++ / py)	98.49	0.014 / 0.225
	Linear Regression	98.38	2.14
	Time Series	98.47	71
k = 1	X-Sketch (C++ / py)	90.71	0.013 / 0.222
	Linear Regression	91.67	2.11
	Time Series	93.67	37.6
k = 2	X-Sketch (C++ / py)	85.31	0.015 / 0.228
	Linear Regression	85.67	2.13
	Time Series	95.36	71.8



06 Conclusion

- 1. k-simplex items
- 2. X-Sketch

Key technique: Short-Term Filtering & Weight Election

- 3. Theoretical & experimental results
- 4. Accelerating machine learning



Source code: https://github.com/SimpleX-Sketch/X-Sketch

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